Some ideas and open problems in real-time stochastic scheduling

Liliana CUCU, TRIO team, Nancy, France

Real-time systems



- Correct reaction
- **I**emporal constraints











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	Real	-tim	ne m	odel	•	$\tau_i = (O_i, C_i, T_i, D_i)$					
$ au_1$	=(C	D_1, C	$T_1, T_1,$	$D_1)$ =	=(1,	2,5,4	4)		↑ ↓	release ti deadline	imes es
-1											
	0		2	3	4	5	6			temps	5





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Why stochastic?

Soft real-time constraints Uncertainness Worst-case behavior is a rare event

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Extracting quantitative information, i.e., obtaining distribution functions

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Temporal analysis of systems with at least one parameter given by a random variables

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Extracting quantitative information



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Joint work with N. Navet and René Schott (TRIO, Nancy)

Activation model of tasks not known

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Monte-Carlo simulation

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Monte-Carlo simulation
Analytical approaches
Markov's, Tchebychev's, Chemoff's upper bounds
Large deviation

- better suited than simulation to rare events
- easily implementable
- embedded in a broader analysis

Large deviation : main result

 $M_n = \frac{1}{n} \sum_{k=1}^n R_{i,k}$ mean of response times over n task instances

$$P(M_n \ge value)$$

Cramer's theorem : if $R_{i,n}$ independent identically distributed random variables

$$P(M_n \in \mathbb{G}) \asymp e^{-n \inf_{x \in \mathbb{G}} I(x)} \mathbb{G} = [value, \infty)$$

$$I(x) = \sup_{\tau > 0} [\tau x - \log E(e^{\tau x})] = \sup_{\tau > 0} [\tau x - \log \sum_{k = -\infty}^{+\infty} p_k e^{k\tau}]$$

Technical contribution

Can deal with distributions given as histograms



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12/26

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Satisfied deadline $satisfyDeadline_i = \begin{pmatrix} yes & no \\ 0.8 & 0.2 \end{pmatrix}$

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$$J_i = \begin{pmatrix} 2 & 3 & 4 \\ 0.7 & 0.1 & 0.3 \end{pmatrix}$$

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Joint work with E. Tovar (Hurray, Portugal)



When minimal inter-arrival times are considered

$$R_i = C_i + \sum_{j \in hp(i)} \left\lceil \frac{R_i}{T_j} \right\rceil C_j$$

This time ...

Response time

$$\mathbb{R}_i = C_i \otimes (\otimes_{k \in P} \lceil \frac{\mathbb{R}_i}{C_k}) \otimes (\otimes_{k \in R} N_{\tau_k} C_k)$$

Algorithm providing a solution

Initial value
$$\Re_i^0 = \begin{pmatrix} 0 \\ 1 \end{pmatrix}$$
 $N_{\tau_k} = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \forall k \in P_{hp(i)}$

New arrivals are added Response time contains to the response time only changed values

NO

YES

All values unchaged or deadline missed ?

 $\begin{pmatrix} 6 & 9 & 11 \\ 0.5 & 0.3 & 0.2 \end{pmatrix}$

 $\mathbb{R}_i =$

Initial values

$$\Re_i^0 = \begin{pmatrix} r_{i,1}^0 \\ 1 \end{pmatrix} = \begin{pmatrix} 0 \\ 1 \end{pmatrix} \text{ and } N_k = \begin{pmatrix} 0 \\ 1 \end{pmatrix}, \forall k \in R_{hp(i)}$$

Iteration *m* - first step

Working random variable
$$L^m$$

 $L_j^m = C_i + \sum_{k \in P_{hp(i)}} \left[\frac{L_j^m}{T_k} \right] \cdot C_k + \sum_{k \in R_{hp(i)}} N_k(r_j^{m-1}) \cdot C_k$

 r_j^{m-1} initial value

An example



$$L_{4}^{1} = \begin{pmatrix} l_{1}^{1} \\ 1 \end{pmatrix} = \begin{pmatrix} 5 \\ 1 \end{pmatrix} \text{ where } l_{1}^{1} \text{ solution of equation } l_{1}^{1} = 0 \cdot C_{1} + \left[\frac{l_{1}^{1}}{T_{2}} \right] C_{2} + 0 \cdot C_{3} + 1 \cdot C_{4}$$

with $r_{1}^{0} = 0$ initial value

Iteration *m* - second step

$\Re_i^m = L^m \otimes (\bigotimes_{k \in R_{hp(i)}} \Delta_k \cdot C_k)$

Back to the example

Task	Т	С
$oldsymbol{ au}_1$	$T_1 = \begin{pmatrix} 8 & 10 & 15 \\ 0.1 & 0.3 & 0.6 \end{pmatrix}$	3
$ au_2$	$T_2 = \begin{pmatrix} 10\\1 \end{pmatrix}$	3
$ au_3$	$T_3 = \begin{pmatrix} 15 & 20\\ 0.6 & 0.4 \end{pmatrix}$	2
${oldsymbol{ au}}_4$	$T_4 = \begin{pmatrix} 15\\1 \end{pmatrix}$	2
$ au_{5}$	$T_5 = \begin{pmatrix} 14 & 22 \\ 0.4 & 0.6 \end{pmatrix}$	2

$$\mathfrak{R}_4^2 = L^2 \otimes \begin{pmatrix} 1 & 2 \\ 0.60.4 \end{pmatrix} C_1 \otimes \begin{pmatrix} 0 \\ 1 \end{pmatrix} C_3$$

Iteration *m* - get ride of unchanged values

$$\begin{cases} L^{m} = \begin{pmatrix} 1 & 3 & 4 \\ 0.50.20.3 \end{pmatrix} \\ \Re_{i}^{m} = \begin{pmatrix} 1 & 2 & 3 & 4 & 5 \\ 0.10.40.30.10.1 \end{pmatrix} \end{cases}$$

New $\Re_{i}^{m} = Comp(\Re_{i}^{m}, L^{m}) = \begin{pmatrix} 2 & 5 \\ 0.4 & 0.1 \end{pmatrix}$

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One entire iteration (3) of our example

The periodic higher tasks are giving a response time:

$$\mathbf{I} = \begin{pmatrix} 20\\ 0.5 \end{pmatrix}$$

The random higher tasks are giving a response time:

$$\Re^{3}_{n,0} = I \otimes \left(F^{*}(20) \cdot C_{3} \right) = \begin{pmatrix} 20 & 21 \\ 0.42 & 0.08 \end{pmatrix},$$

where $F^{*}(20) = \begin{pmatrix} 2 & 3 \\ 0.84 & 0.16 \end{pmatrix}$

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 Next step?

Marcolition: deterministic case

Initial condition: deterministic case Robustness condition: worst-case behavior of the algorithms is rare

Initial condition: deterministic case
 Robustness condition: worst-case
 behavior of the algorithms is rare
 Worst-case condition: insure worst-case
 when mixing hard and soft

A. Burns, G. Bernat, I. Broster - A probabilistic framework for schedulability analysis

M. Pinedo

S. Edgar, A. Burns N. Navet, L. Cucu, R. Schott

J.L. Diaz, D.F. Garcìa, K. Kim, C.-G. Lee, L. Lo Bello, J.M. Lopez, S.L. Min, O. Mirabella L. Cucu, E.Tovar

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- Statistical analysis of WCET for scheduling
- Probabilistic estimation of response times through large deviations
- Stochastic analysis of real-time systems
 - A framework for response times analysis of
- fixed-priority tasks with stochastic inter-arrival times

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Thank you for your attention



Open problems in stochastic real-time scheduling : introduction of dependent random variables